COMP312-09A Communications and Systems Software

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COMP312 - ADSL

Motivation

- Sometimes we want to tunnel one protocol over another protocol
 - Maybe the network does not understand how to forward that protocol
 - Maybe want to route around a failure
 - Maybe want to tunnel for policy reasons
 - Mobility
 - VPN services



VPN

- · Virtual Private Network
- Point-to-point network link routed over existing available networks.
 - Cheaper in terms of capital expenditure
 - Easy to add

VPNs and ISPs

- Useful for whole sale reselling ADSL services
- Tunnel all clients of an ISP across the wholesale core network to the ISP
- · Clients appear to be directly connected to the ISP



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Tunnelling protocols

- There are many tunneling protocols
- · In this lecture
 - GRE
 - L2TP



Generic Routing Encapsulation

- GRE
- Defined in RFCs 1701, 2784, 2890
- Encapsulated in IP, protocol number 47
- · Generic protocol for tunnelling
- Though only IPv4 tunnelling is well defined and widely implemented



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0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 1 2 3 4 5 6 7 8 9 0 1 2 2 3 4 5 6 7 8 9 0

GRE Header: RFC 1701

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GRE

RFC 1701

- · C: checksum bit
- R: routing bit
- · K: key bit
- · S: sequence bit
- s: strict source route
- · Recur: recursion control
- · Protocol: Ethernet protocol type field
 - 0x0800 for IPv4



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GRE Header: RFC 2784

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GRE in Action

- GRE is normally stateless
 - No information about remote endpoint status
 - Up if local interface is configured and remote IP is routeable
- Some GRE implementations support Keepalives
 - No need for other end support
 - Two GRE Headers, one forward and one reverse
 - Other end decapsulates forward header and forwards based on reverse

PPTP

- Point-to-point tunnelling protocol
- Informational RFC2637
 - Vendor written (MS, 3COM and others)
 - PPP encapsulated over IP
 - PPP negotiation negotiated over TCP port 1723
 - Tunnelled packets use GRE
 - GRE version field set to one, GRE header slightly modified.



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GRE Issues - RFC 2784

- · Path MTU Discovery
- IPv6
- · Interaction with ICMP
- Looping



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PPTP Issues

- Multiple connections
- Security



Motivation - GRE limitations

- ISPs require authentication for billing
 - GRE authentication not standardised, removed from later versions
- Interaction between other L3 protocols and GRE not well defined
- · Path-MTU discovery problems

Motivation

- It would be easier to just tunnel PPP sessions over IP
 - IP is ubiquitous
 - PPP has well defined authentication methods
 - PPP defines NCPs for tunnelling various protocols
 - PPP has a large installed base



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L2TP

- · Layer-2 Tunnelling Protocol
 - L2TPv2 defines tunnelling PPP circuits
 - RFC 2661
 - L2TPv3 defines tunnelling other L2 protocols
 - · "Pseudo-wire"
 - · Ethernet, Frame relay
 - RFC 3931



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LAC: L2TP Access Concentrator LNS: L2TP Network Server LAC LNS

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L2TP protocol overview

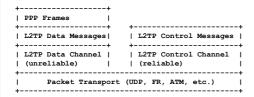


Figure 3.0: L2TP Protocol Structure, RFC 2661



L2TP Control Connection

- · Tunnel setup and maintenance
- Uses UDP, implements TCP-like windowing and congestion control
- Three message exchange
 - $\ SCCRQ: \ Start-Control-Connection-Request$
 - SCCRP: Start-Control-Connection-Reply
 - SCCCN: Start-Control-Connection-Connected

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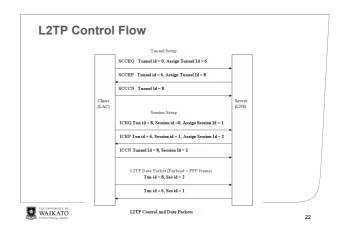
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L2TP Call Control

- · Tunnel individual PPP sessions
- · Unreliable UDP stream of packets
- · Three-packet exchange
 - ICRQ: Incoming-Call-Request - ICRP: Incoming-Call-Reply - ICCN: Incoming-Call-Connected





L2TP message format

- 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9
 - T: Type S: Sequence P: Priority L: Lenath O: Offset Ver: Version
 - Identifier of control connection Identifer for session within tunnel Tunnel ID: Session ID: Ns: Nr:

expected next sequence number expected next sequence number for control messages





- 2 Options
- · Automatic forwarding
 - PPP session established to LAC.
 - Forwarded automatically to LNS based on configuration or other (e.g. dialling) info.
 - LNS Authenticates
- · Partial authentication
 - PPP session established to LAC
 - LAC starts authentication
 - Forwarded to LNS based on authentication details
 - Authentication details forwarded and LNS checks



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L2TP and MTU

- L2TP, like other tunnelling mechanisms, introduces per-packet framing overhead
 - Reduces MTU
- · Two solutions
 - PPP negotiation of MTU
 - Provision network between LAC and LNS to encapsulate commonly required packet size
 - 1500 user bytes + PPP header + L2TP + UDP + IP



L2TP: LAC to LNS security

- Authentication
 - CHAP using shared secret
- Encryption
 - IPSec
 - User-PPP + L2TP + UDP + IPSec + IP



L2TPv3

- Generic version of L2TP
 - Can be used for tunneling any protocol, not just PPP
- RFC3931 is base specification
- Separate RFCs for encapsulation types
 - PPP, Ethernet, Frame Relay etc
- 32 bit tunnel ID to support more sessions
- Improved Authentication

