


COMP340-08B Reasoning About Programs

8. Natural Deduction

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Rules for \wedge

\wedge -Elimination

(\wedge -elim-1) $\frac{A \wedge B}{A}$ (\wedge -elim-2) $\frac{A \wedge B}{B}$

\wedge -Introduction

(\wedge -intro) $\frac{A \quad B}{A \wedge B}$

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Natural Deduction

- A formal framework to derive proofs in propositional and predicate logic.
- Produces proofs that are similar in style to common (mathematical) reasoning.
- Sound and complete.
- Today: Only propositional logic.

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Rules for \rightarrow

\rightarrow -Elimination (Modus Ponens)

(\rightarrow -elim) $\frac{A \rightarrow B \quad A}{B}$

\rightarrow -Introduction (Conditional Proof)

(\rightarrow -intro) $\frac{A}{B}$
 $\frac{B}{A \rightarrow B}$

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Elimination and Introduction

Natural deduction uses two types of inference rules for each connective:

- **Elimination rules:**
Remove a connective from a formula.
- **Introduction rules:**
Add a connective to a formula.

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Example Proof

Premises: $p, q, p \wedge q \rightarrow r \wedge s$

Conclusion: s

1.	p	premise
2.	q	premise
3.	$p \wedge q \rightarrow r \wedge s$	premise
4.	$p \wedge q$	\wedge -intro: 1,2
5.	$r \wedge s$	\rightarrow -elim: 4,3
6.	s	\wedge -elim-2: 5

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Remember

- A **proof** is a step-by-step demonstration that the conclusion follows from the premises.
- Each step is
 - numbered (for future references) and
 - justified by an appropriate rule.

Assumed Premises

- Some rules introduce new premises.
- These temporary premises are called **assumed premises** or **assumptions**.
- They are only available in a sub-proof and must be **discharged**, i.e., forgotten when the sub-proof is completed.

Conditional Proof

$$(\rightarrow\text{-intro}) \frac{\frac{A}{B}}{A \rightarrow B}$$

In order to prove $A \rightarrow B$, we assume that A holds and prove B from this assumption.

Rules for \vee

\vee -Introduction

$$(\vee\text{-intro-1}) \frac{A}{A \vee B} \quad (\vee\text{-intro-2}) \frac{B}{A \vee B}$$

\vee -Elimination (Proof by Cases)

$$(\vee\text{-elim}) \frac{A \vee B \quad \frac{A}{C} \quad \frac{B}{C}}{C}$$

Example Using Conditional Proof

- | | | |
|----|--|---------------------------|
| 1. | $p \rightarrow r$ | assumption |
| 2. | $p \wedge q$ | assumption |
| 3. | p | \wedge -elim-1: 2 |
| 4. | r | \rightarrow -elim: 1,3 |
| 5. | $p \wedge q \rightarrow r$ | \rightarrow -intro: 2-4 |
| 6. | $(p \rightarrow r) \rightarrow (p \wedge q \rightarrow r)$ | \rightarrow -intro: 1-5 |

Proof by Cases: an Example

- Premise:** $(p \wedge s) \vee (q \wedge s)$
Conclusion: s
- | | | |
|----|----------------------------------|-------------------------|
| 1. | $(p \wedge s) \vee (q \wedge s)$ | premise |
| 2. | $p \wedge s$ | assumption |
| 3. | s | \wedge -elim-2: 2 |
| 4. | $q \wedge s$ | assumption |
| 5. | s | \wedge -elim-2: 4 |
| 6. | s | \vee -elim: 1,2-3,4-5 |

Rules for \leftrightarrow

\leftrightarrow -Elimination

$$(\leftrightarrow\text{-elim-1}) \frac{A \leftrightarrow B}{A \rightarrow B} \quad (\leftrightarrow\text{-elim-2}) \frac{A \leftrightarrow B}{B \rightarrow A}$$

\leftrightarrow -Introduction

$$(\leftrightarrow\text{-intro}) \frac{A \rightarrow B \quad B \rightarrow A}{A \leftrightarrow B}$$

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Natural Deduction — Tips

- Apply all possible elimination rules to all premises and assumptions
This gives more and more premises that will help you to complete the proof.
- Apply the appropriate introduction rule to the conclusion
This may produce new assumptions or conclusions – then repeat these two steps.

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Example with Equivalence

1.	$p \wedge q$	assumption
2.	p	\wedge -elim-1: 1
3.	q	\wedge -elim-2: 1
4.	$q \wedge p$	\wedge -intro: 2,3
5.	$(p \wedge q) \rightarrow (q \wedge p)$	\rightarrow -intro: 1–4
6.	$q \wedge p$	assumption
7.	q	\wedge -elim-1: 6
8.	p	\wedge -elim-2: 6
9.	$p \wedge q$	\wedge -intro: 8,7
10.	$(q \wedge p) \rightarrow (p \wedge q)$	\rightarrow -intro: 6–9
11.	$(p \wedge q) \leftrightarrow (q \wedge p)$	\leftrightarrow -intro: 5,10

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JAPE – a Proof Assistant

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Proving Equivalences

To establish an equivalence

$$A \leftrightarrow B$$

we have to prove two implications:

- $A \rightarrow B$ – Assume A and prove B
- $B \rightarrow A$ – Assume B and prove A

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Reading

Huth & Ryan:
Section 1.2
pp. 5–31.

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