


**COMP340-08B**  
**Reasoning**  
**About Programs**

**13. Proof Writing**  
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**Common Phrases in Proofs (3)**

**$\forall$ -Introduction**  
 Let  $\langle item \rangle$  be an arbitrary  $\langle element \rangle$ .  
 ...  
 Therefore  $\langle item \rangle$  satisfies  $\langle the claim \rangle$ .  
 Since  $\langle item \rangle$  was chosen arbitrarily, it follows  
 that every  $\langle element \rangle$  satisfies  $\langle the claim \rangle$ .

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**Common Phrases in Proofs (1)**

**$\rightarrow$ -Introduction**  
 Assume that  $\langle precondition \rangle$ .  
 ...  
 Therefore we have  $\langle conclusion \rangle$ .  
 It follows that  $\langle precondition \rangle$  always implies  
 $\langle conclusion \rangle$ .

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**Common Phrases in Proofs (5)**

**$\vee$ -Elimination**  
 (Given some disjunction, say  $\langle case1 \vee case2 \rangle$ .)  
 Consider the following cases.  
 First assume that  $\langle case1 \rangle$  holds.  
 ...  
 Then it follows that  $\langle the claim \rangle$ .  
 Second assume that  $\langle case2 \rangle$  holds.  
 ...  
 Then it follows that  $\langle the claim \rangle$ .  
 In all cases we have  $\langle the claim \rangle$ .

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**Common Phrases in Proofs (2)**

**$\leftrightarrow$ -Introduction**  
 First assume that  $\langle condition1 \rangle$ .  
 ...  
 Therefore we have  $\langle condition2 \rangle$ .  
 Second assume that  $\langle condition2 \rangle$ .  
 ...  
 Therefore we have  $\langle condition1 \rangle$ .  
 It follows that  $\langle condition1 \rangle$  and  $\langle condition2 \rangle$  are  
 equivalent.

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**Common Phrases in Proofs (6)**

***false*-Elimination,  $\neg$ -Introduction**  
 Assume that  $\langle the claim \rangle$  does not hold.  
 ...  
 This is a contradiction.  
 Therefore, the assumption was wrong,  
 and it follows that  $\langle the claim \rangle$  holds.

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### Another Example: Semigroups

A **semigroup** is an algebraic structure that involves an associative operation.

Some set

}

Associative binary operation

$(G, *)$

**Axiom**

$$\forall x \forall y \forall z \quad x * (y * z) = (x * y) * z$$

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### Groups

A semigroup  $(G, *)$  is called a **group** if

- $G$  contains a left identity, and
- every element of  $G$  has a left inverse.

**Known Properties of groups:**

- A group has exactly one identity element, which is both a left and a right identity.
- Every element of a group has exactly one inverse, which is both its left and right inverse.

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### Examples of Semigroups

**Examples**

- $(\mathbb{N}, +)$  – Natural numbers with addition
- $(\mathbb{R}, \cdot)$  – Real numbers with multiplication
- Functions  $f: \mathbb{N} \rightarrow \mathbb{N}$  with functional composition.

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### A Proof in Group Theory

**Proposition.**  
Let  $(G, *)$  be a semigroup. Then every right inverse of an element of  $G$  is equal to every left inverse of that element.

**Proof.**  
Let  $a \in G$ , let  $l$  be a left inverse of  $a$ , and let  $r$  be right inverse of  $a$ . Then  $l * a$  is a left identity, and  $a * r$  is a right identity. It follows that  $l = l * (a * r) = (l * a) * r = r$ , i.e.,  $l = r$ . Since  $a, l$ , and  $r$  were chosen arbitrarily, it follows that every right inverse is equal to every left inverse. □

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### Identity and Inverse Elements

Let  $(G, *)$  be a semigroup and  $e \in G$ .

- $e$  is a **left identity**, if  $\forall x \quad e * x = x$ .
- $e$  is a **right identity**, if  $\forall x \quad x * e = x$ .
- $e$  is an **identity**, if  $e$  is a right and a left identity.

Let  $a \in G$ .

- $a'$  is a **left inverse** of  $a$  if  $a' * a$  is a left identity.
- $a'$  is a **right inverse** of  $a$  if  $a * a'$  is a right identity.
- $a'$  is an **inverse** of  $a$  if  $a'$  is a both right and left inverse of  $a$ .

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### The Proof in Natural Deduction

1. $\forall x \forall y \forall z \quad (x * y) * z = x * (y * z)$	Premise
2. $\forall x \quad (\text{left\_id}(x) \rightarrow \forall y \quad x * y = y)$	Premise
3. $\forall x \quad (\text{right\_id}(x) \rightarrow \forall y \quad y * x = x)$	Premise
4. $\forall x \forall y \quad (\text{left\_inv}(x, y) \rightarrow \text{left\_id}(x * y))$	Premise
5. $\forall x \forall y \quad (\text{right\_inv}(x, y) \rightarrow \text{right\_id}(y * x))$	Premise
...	
n. $\forall x \forall y \forall z \quad (\text{left\_inv}(y, x) \wedge \text{right\_inv}(z, x) \rightarrow y = z)$	

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