


# COMP340-08B Reasoning About Programs

## 14. Proof by Induction

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### Principle of Induction

$$(p(0) \wedge \forall x (\text{nat}(x) \wedge p(x) \rightarrow p(s(x)))) \rightarrow \forall x (\text{nat}(x) \rightarrow p(x))$$

To prove a property **P** of all natural numbers:

**Inductive base.** Prove that 0 satisfies **P**.

**Inductive hypothesis.** Assume that some number  $n$  satisfies **P**.

**Inductive step.** Using the inductive hypothesis, prove that  $n + 1$  satisfies **P**.

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### About the RISC Proof Navigator

- Type ProofNavigator to start
- Homepage and download  
<http://www.risc-uni-linz.ac.at/research/formal/software/ProofNavigator>
- Local copies of documentation on the COMP340-08B homepage in Moodle:
  - Online manual available from Help menu or from the homepage
  - Read the tutorial (section 3.1 of the manual)

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### Example

Proving  
**(Claim)**  $\forall x:\mathbb{N} \ x + 0 = x$   
from the axioms  
**(A1)**  $\forall x:\mathbb{N} \ 0 + x = x$   
**(A2)**  $\forall x:\mathbb{N}, y:\mathbb{N} \ s(x) + y = s(x + y)$

**Inductive base.**  
 $0 + 0 = 0$  by axiom (A1) By Axiom (A2)

**Inductive hypothesis.**  
 $n + 0 = n$  for some  $n \in \mathbb{N}$  By Inductive hypothesis

**Inductive step.**  
 $s(n) + 0 = s(n + 0) = s(n)$  □

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### Example Proofs with ProofNavigator

**Definition**  
 $\text{divides}(x, y) \equiv \exists n:\mathbb{N} \ n \cdot x = y$

**Theorems to prove**

**(DIV<sub>0</sub>)**  $\forall x:\mathbb{N} \ \text{divides}(x, 0)$

**(DIV<sub>self</sub>)**  $\forall x:\mathbb{N} \ \text{divides}(x, x)$

**(DIV<sub>1</sub>)**  $\forall x:\mathbb{N} \ \text{divides}(1, x)$

```
expand divides;
scatter;
instantiate 0 in h36;
instantiate x_0 in gwo;
```

```
scatter;
expand divides;
instantiate 1 in gpq;
lemma MULT_la;
auto xdm;
```

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### Reading on Induction

Huth & Ryan:  
Section 1.4.2  
pp. 40–45.

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### To be Proven by Induction

$\forall x:\mathbb{N} \ x + 0 = x$   
 $\forall x:\mathbb{N}, y:\mathbb{N}, z:\mathbb{N} \ (x + y) + z = x + (y + z)$   
 $\forall x:\mathbb{N}, y:\mathbb{N} \ x + y = y + x$   
 $\forall x:\mathbb{N} \ x \cdot 0 = 0$   
 $\forall x:\mathbb{N} \ x \cdot 1 = x$   
 $\forall x:\mathbb{N}, y:\mathbb{N}, z:\mathbb{N} \ (x + y) \cdot z = x \cdot z + y \cdot z$   
 $\forall x:\mathbb{N}, y:\mathbb{N}, z:\mathbb{N} \ x \cdot (y + z) = x \cdot y + x \cdot z$   
 $\forall x:\mathbb{N}, y:\mathbb{N}, z:\mathbb{N} \ (x \cdot y) \cdot z = x \cdot (y \cdot z)$   
 $\forall x:\mathbb{N}, y:\mathbb{N} \ x \cdot y = y \cdot x$

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### ProofNavigator Proof Script

```

induction x in cql;
scatter;
auto;
scatter;
instantiate x_0, mult(y_0,z_0) in fcb;
instantiate y_0, z_0 in bwe;
lemma MULT_dist;
instantiate mult(x_0,y_0), y_0, z_0 in x3a;
instantiate x_0, y_0 in fcb;
    
```

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### Another Example

Proving

**(Claim)**  $\forall x:\mathbb{N}, y:\mathbb{N}, z:\mathbb{N} \ x \cdot (y \cdot z) = (x \cdot y) \cdot z$   
from the axioms

**(M1)**  $\forall x:\mathbb{N} \ 0 \cdot x = 0$

**(M2)**  $\forall x:\mathbb{N}, y:\mathbb{N} \ (x + 1) \cdot y = x \cdot y + y$

by induction on  $x \dots$

**Inductive base.**

$0 \cdot (y \cdot z) = 0 = 0 \cdot z = (0 \cdot y) \cdot z$  by axiom **(M1)**

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### Yet Another Example

Proving

**(Claim)**  $\forall x:\mathbb{N}, y:\mathbb{N} \ x \cdot y = y \cdot x$

by induction on  $x \dots$

**Inductive base:**  $x = 0$

$0 \cdot x = 0 = x \cdot 0$  using lemma  $\forall x:\mathbb{N} \ x \cdot 0 = 0$

**Inductive hypothesis:**  $x = n$

$\forall y:\mathbb{N} \ n \cdot y = y \cdot n$  for some  $n \in \mathbb{N}$

**Inductive step:**  $x = n \rightsquigarrow x = n + 1$

is proven by induction on  $y \dots$

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### Example continued

**Inductive hypothesis.**

$\forall y:\mathbb{N}, z:\mathbb{N} \ n \cdot (y \cdot z) = (n \cdot y) \cdot z$  for some  $n \in \mathbb{N}$

**Inductive step.**

$(n + 1) \cdot (y \cdot z)$   
 $= n \cdot (y \cdot z) + y \cdot z$  by axiom **(M2)**  
 $= (n \cdot y) \cdot z + y \cdot z$  by inductive hypothesis  
 $= (n \cdot y + y) \cdot z$  by distributivity  
 $= ((n + 1) \cdot y) \cdot z$  by axiom **(M2)** □

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### Nested Induction

**Inductive base:**  $y = 0$

$(n + 1) \cdot 0 = 0 = 0 \cdot (n + 1)$  using lemma  $\forall x:\mathbb{N} \ x \cdot 0 = 0$

**Inductive hypothesis:**  $y = m$

$(n + 1) \cdot m = m \cdot (n + 1)$  for some  $m \in \mathbb{N}$

**Inductive step:**  $y = m \rightsquigarrow y = m + 1$

$(n + 1) \cdot (m + 1)$   
 $= n \cdot (m + 1) + (m + 1)$  by axiom **(M2)**  
 $= (m + 1) \cdot n + m + 1$  by inductive hypothesis for  $n$   
 $= m \cdot n + n + m + 1$  by axiom **(M2)**  
 $= n \cdot m + m + n + 1$  by inductive hypothesis for  $n$   
 $= (n + 1) \cdot m + n + 1$  by axiom **(M2)**  
 $= m \cdot (n + 1) + n + 1$  by inductive hypothesis for  $m$   
 $= (m + 1) \cdot (n + 1)$  by axiom **(M2)** □

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